

# Streaming max-min Filter

using no more than 3

comparisons per element

# Preface

在图像去雾算法中，Dr. Kaiming He 给出了 dark channel 的方法，其中第一步是做暗通道提取。本文介绍了一种取最大值最小值的算法

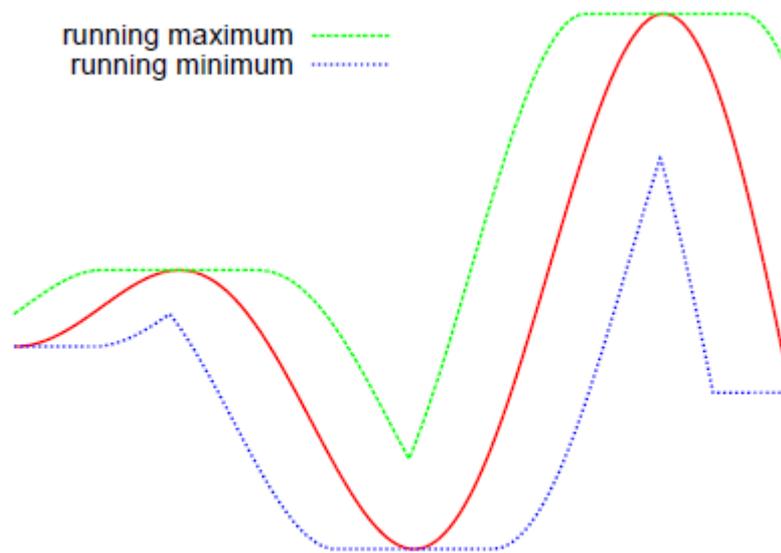
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# Preface

- $N$  个数据需要  $n-1$  次比较得到最大或最小值。
- 需要  $1.5*n-2$  次比较才能得到最大和最小值。

## Running max-min filter:

given an array  $a_1, \dots, a_n$ , find the maximum and the minimum over all windows of size  $w$ , that is  $\max / \min_{i \in [j, j+w)} a_i$  for all  $j$ .



**Fig. 1.1:** Example of a running MAX-MIN filter.

# Running Max-Min Filter

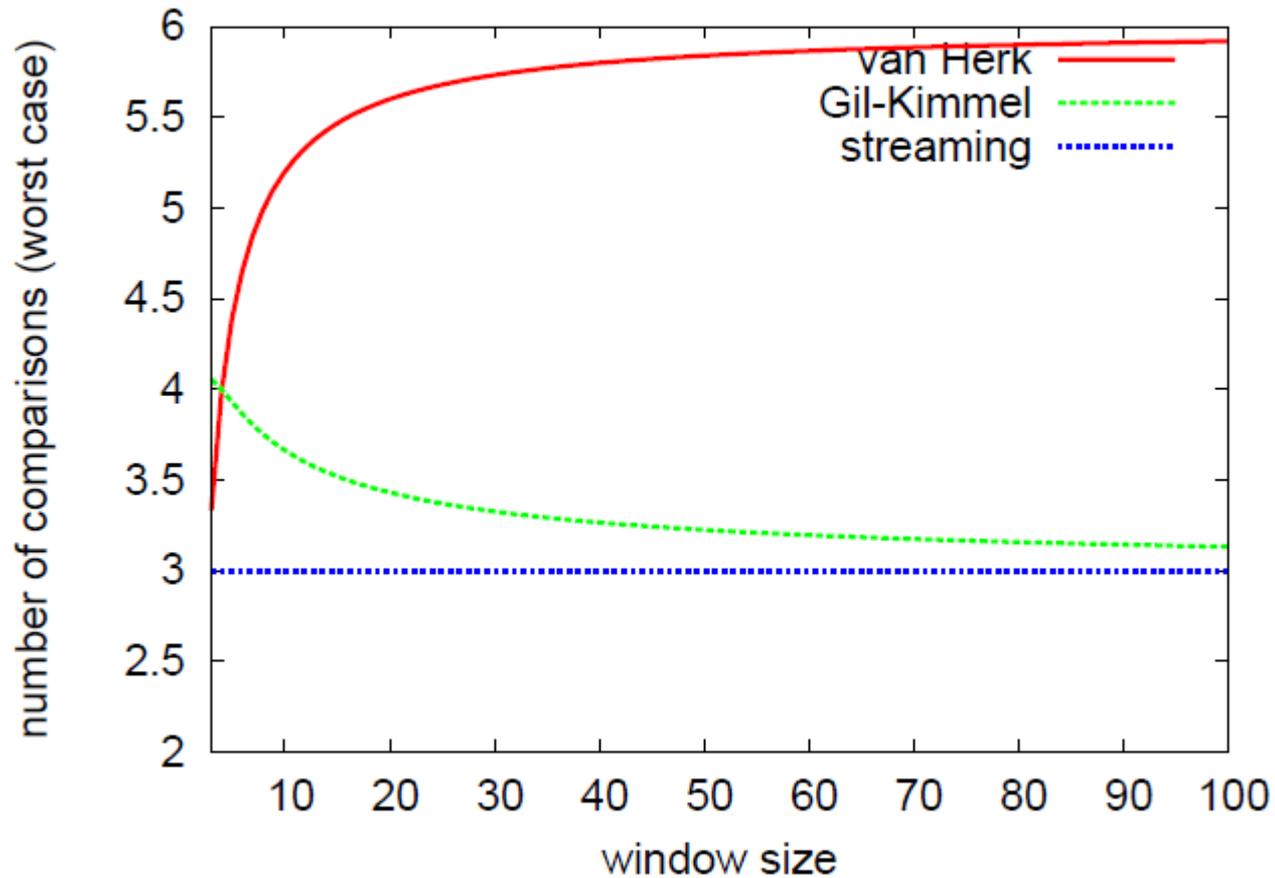
TABLE I: Worst-case number of comparisons and stream latency for competitive MAX-MIN filter algorithms. Stream latency and memory usage (buffer) are given in number of elements.

algorithm	comparisons per element (worst case)	stream latency	buffer
naive	$2w - 2$	0	$O(1)$
van Herk [1992], Gil and Werman [1993]	$6 - 8/w$	$w$	$4w + O(1)$
Gil and Kimmel [2002]	$3 + 2 \log w/w + O(1/w)$	$w$	$6w + O(1)$
<b>New algorithm</b>	<b>3</b>	<b>0</b>	$2w + O(1)$

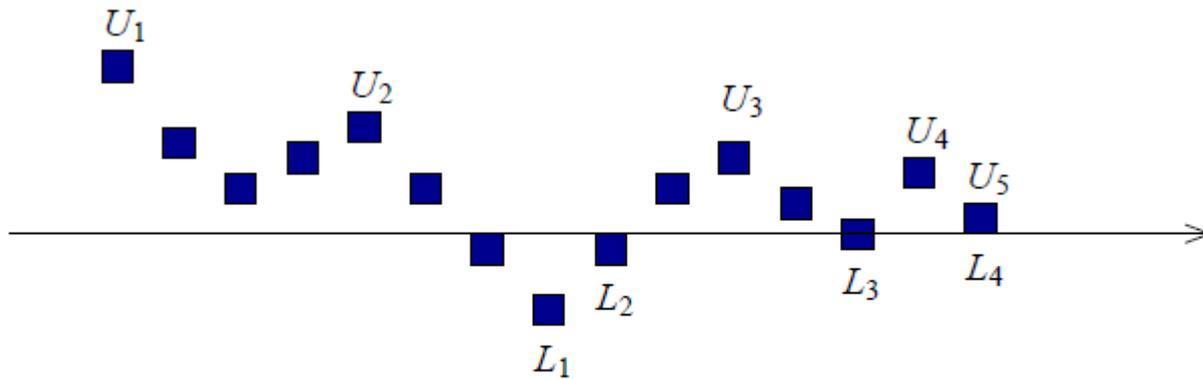
\*当数据流是分段单调的时候，该算法的CPE(worst)可以降到2。

\*stream latency是指延时流。

# Running Max-Min Filter

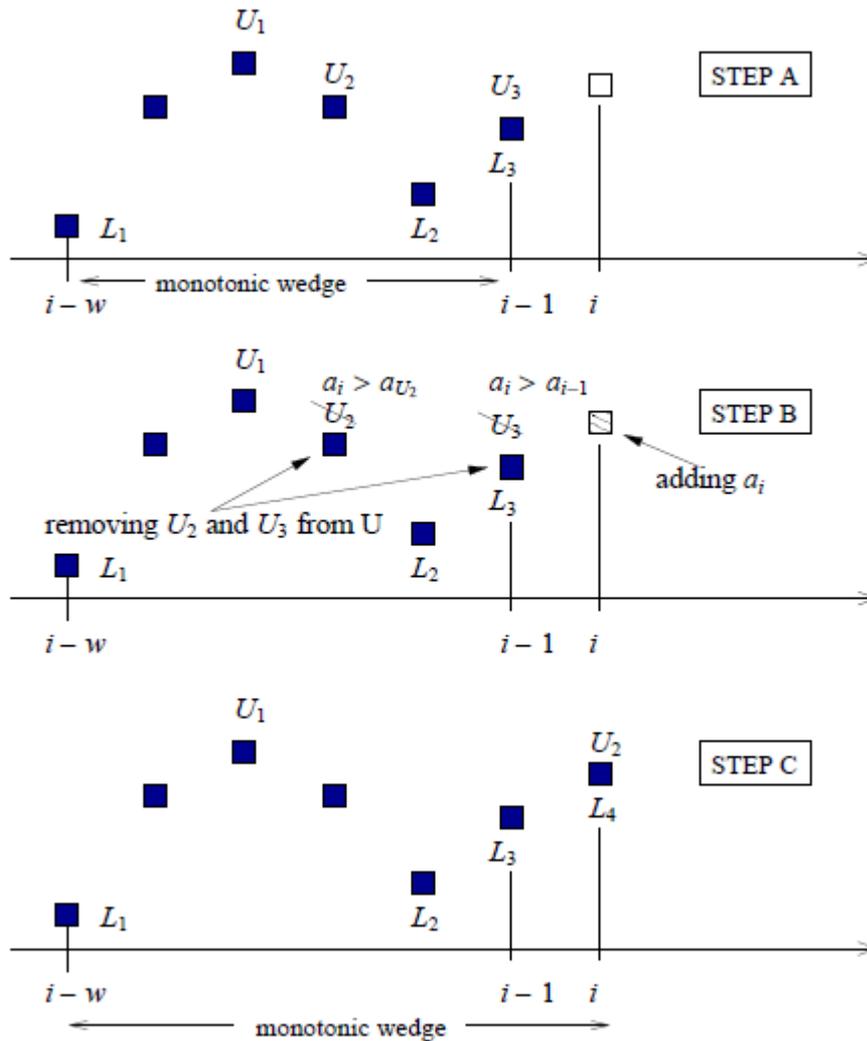


# Running Max-Min Filter



**Fig. 4.3:** Example of a monotonic wedge: data points run from left to right.

# Running Max-Min Filter



# Running Max-Min Filter

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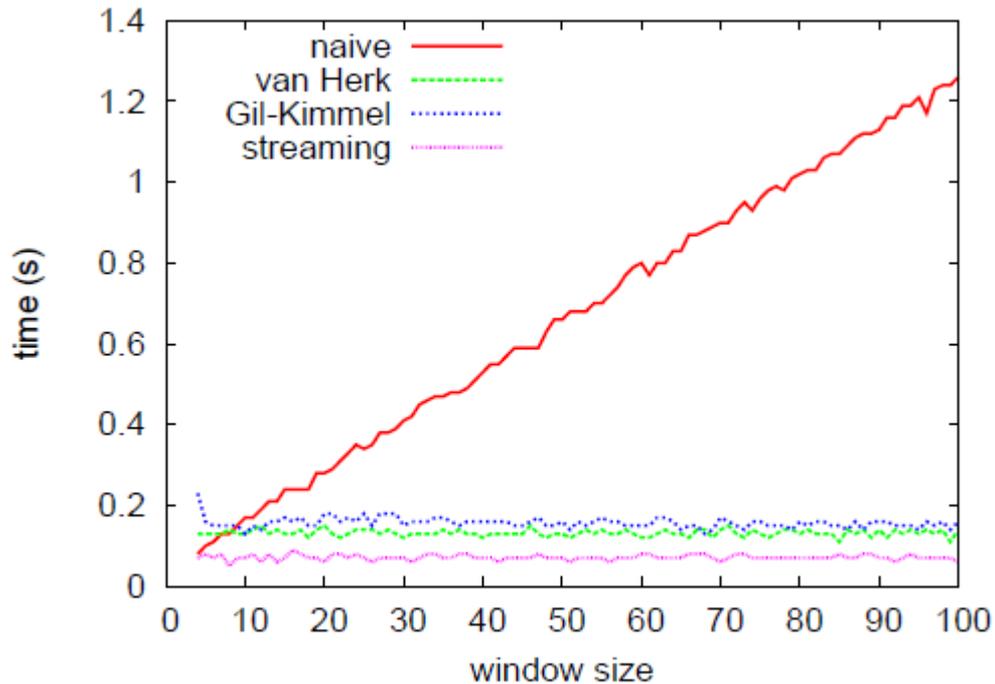
**Algorithm 1** Streaming algorithm to compute the MAX-MIN filter using no more than 3 comparisons per element.

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```
1: INPUT: an array  $a$  indexed from 1 to  $n$ 
2: INPUT: window width  $w > 2$ 
3:  $U, L \leftarrow$  empty double-ended queues, we append to “back”
4: append 1 to  $U$  and  $L$ 
5: for  $i$  in  $\{2, \dots, n\}$  do
6:   if  $i \geq w + 1$  then
7:     OUTPUT:  $a_{\text{front}(U)}$  as maximum of range  $[i - w, i)$ 
8:     OUTPUT:  $a_{\text{front}(L)}$  as minimum of range  $[i - w, i)$ 
9:     if  $a_i > a_{i-1}$  then
10:      pop  $U$  from back
11:      while  $a_i > a_{\text{back}(U)}$  do
12:        pop  $U$  from back
13:     else
14:       pop  $L$  from back
15:       while  $a_i < a_{\text{back}(L)}$  do
16:        pop  $L$  from back
17:     append  $i$  to  $U$  and  $L$ 
18:     if  $i = w + \text{front}(U)$  then
19:       pop  $U$  from front
20:     else if  $i = w + \text{front}(L)$  then
21:       pop  $L$  from front
```

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# Running Max-Min Filter



(a) Input data is a sine wave with a period of 10 000 data points.

# Appendix

```
// input: array a, integer window width w
// output: arrays maxval and minval
// buffer: lists U and L
// requires: STL for deque support
deque<int> U, L;
for(uint i = 1; i < a.size(); ++i) {
    if(i>=w) {
        maxval[i-w] = a[U.size()>0 ? U.front() : i-1];
        minval[i-w] = a[L.size()>0 ? L.front() : i-1];
    } // end if
    if(a[i] > a[i-1]) {
        L.push_back(i-1);
        if(i == w+L.front()) L.pop_front();
        while(U.size()>0) {
            if(a[i]<=a[U.back()]) {
                if(i == w+U.front()) U.pop_front();
                break;
            } // end if
            U.pop_back();
        } // end while
    } else {
        U.push_back(i-1);
        if(i == w+U.front()) U.pop_front();
        while(L.size()>0) {
            if(a[i]>=a[L.back()]) {
                if(i == w+L.front()) L.pop_front();
                break;
            } //end if
            L.pop_back();
        } //end while
    } // end if else
} // end for
maxval[a.size()-w] = a[U.size()>0 ? U.front() : a.size()-1];
minval[a.size()-w] = a[L.size()>0 ? L.front() : a.size()-1];
```